Automated Black-box Verification of Networking Systems



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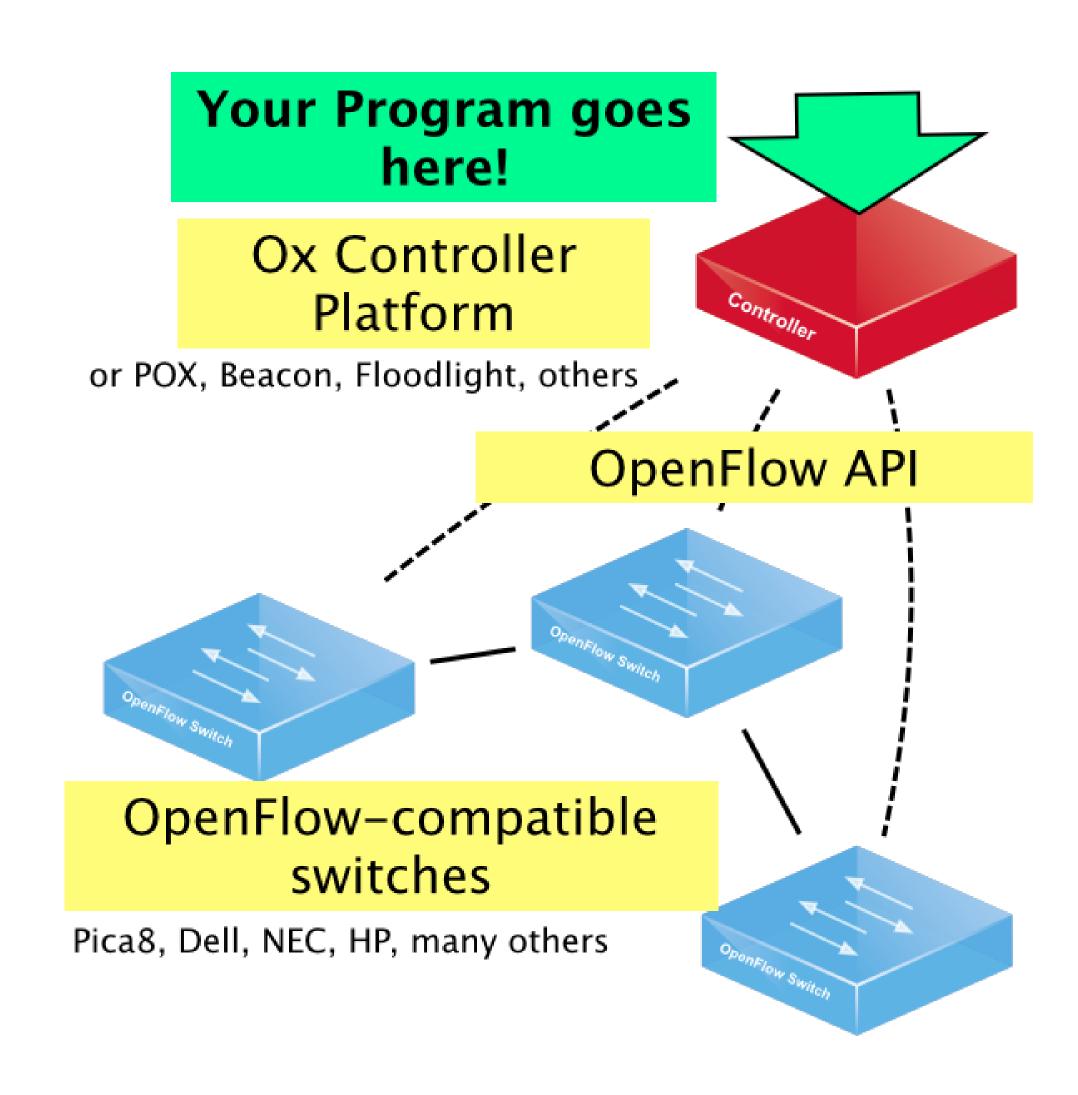


Stefan Zetzsche





Software-Defined Networks



Verification of networks

Trend in PL&Verification after Software-Defined Networks

- Design high-level languages that model essential network features
- Develop semantics that enables reasoning precisely about behaviour
- Build tools to synthesise low-level implementations automatically

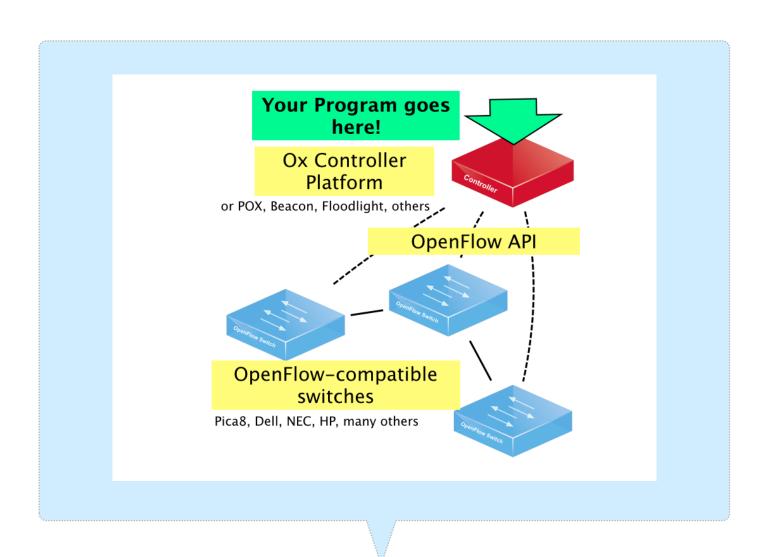
- Frenetic [Foster & al., ICFP 11]
- Pyretic [Monsanto & al., NSDI 13]
- Maple [Voellmy & al., SIGCOMM 13]
- FlowLog [Nelson & al., NSDI 14]
- Header Space Analysis [Kazemian & al., NSDI 12]
- VeriFlow [Khurshid & al., NSDI 13]
- NetKAT [Anderson & al., POPL 14]
- and many others . . .

This is all very nice but..

What if there is no formal model?

Does the low-level implementation really do what it is supposed to do?

What we propose



Build black-box model via interactions with the system

Automated Modelling

Automated Verification

Properties

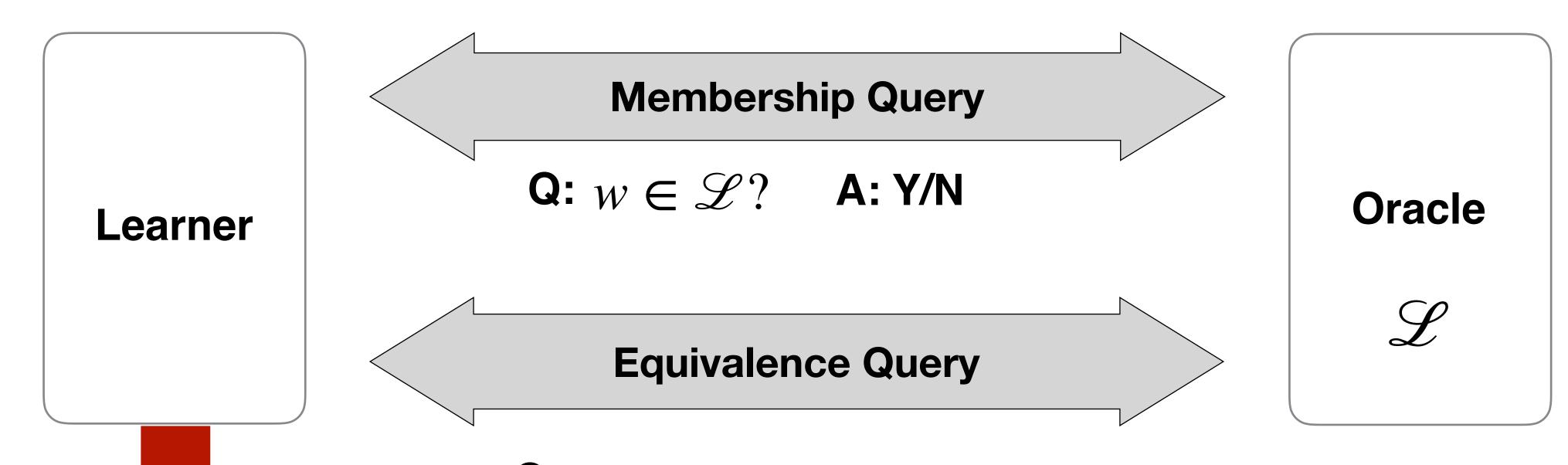
Automata learning (Angluin '87)

Finite alphabet of system's actions A

Minimal DFA

accepting \mathscr{L}

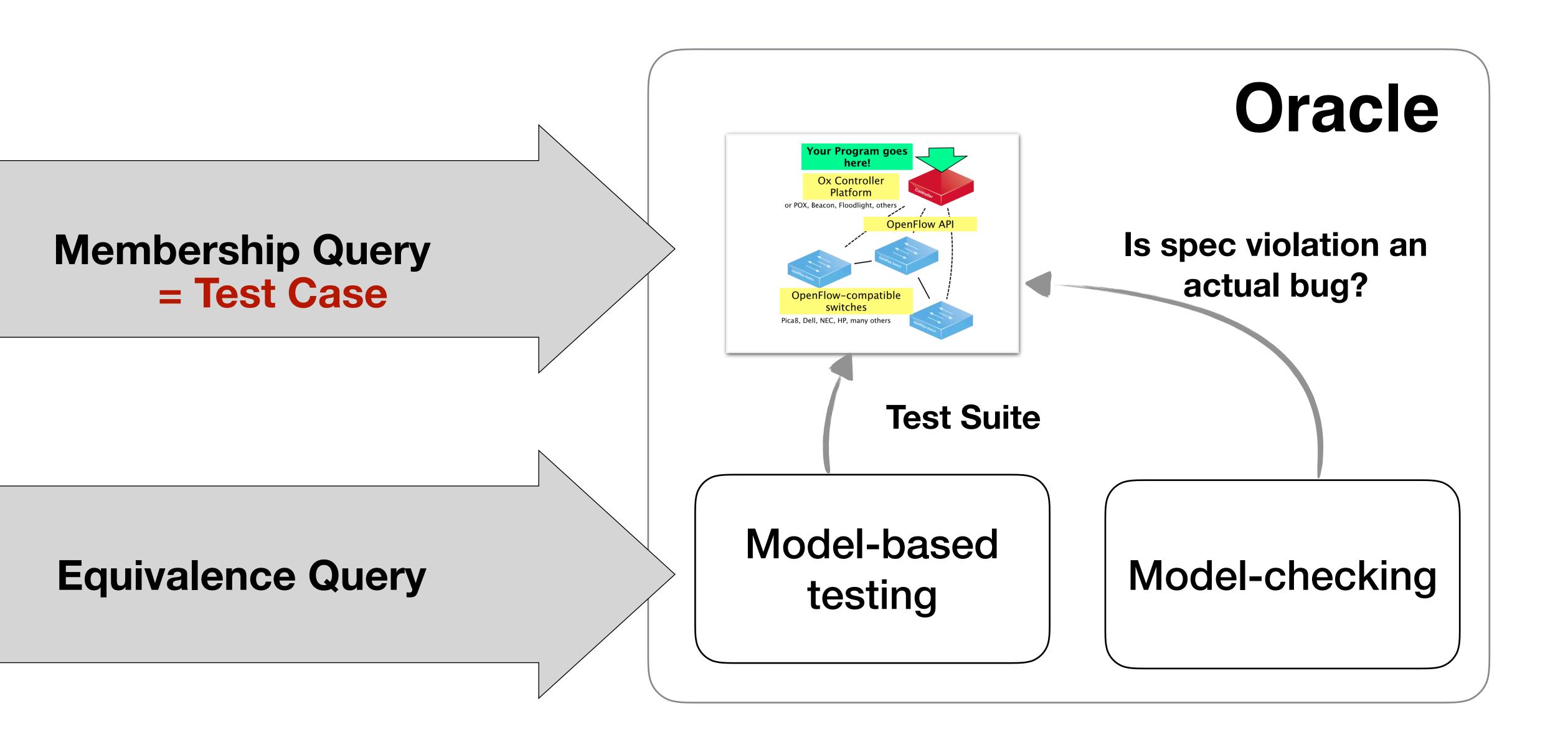
Set of system behaviours is a **regular language** $\mathscr{L} \subseteq A^*$



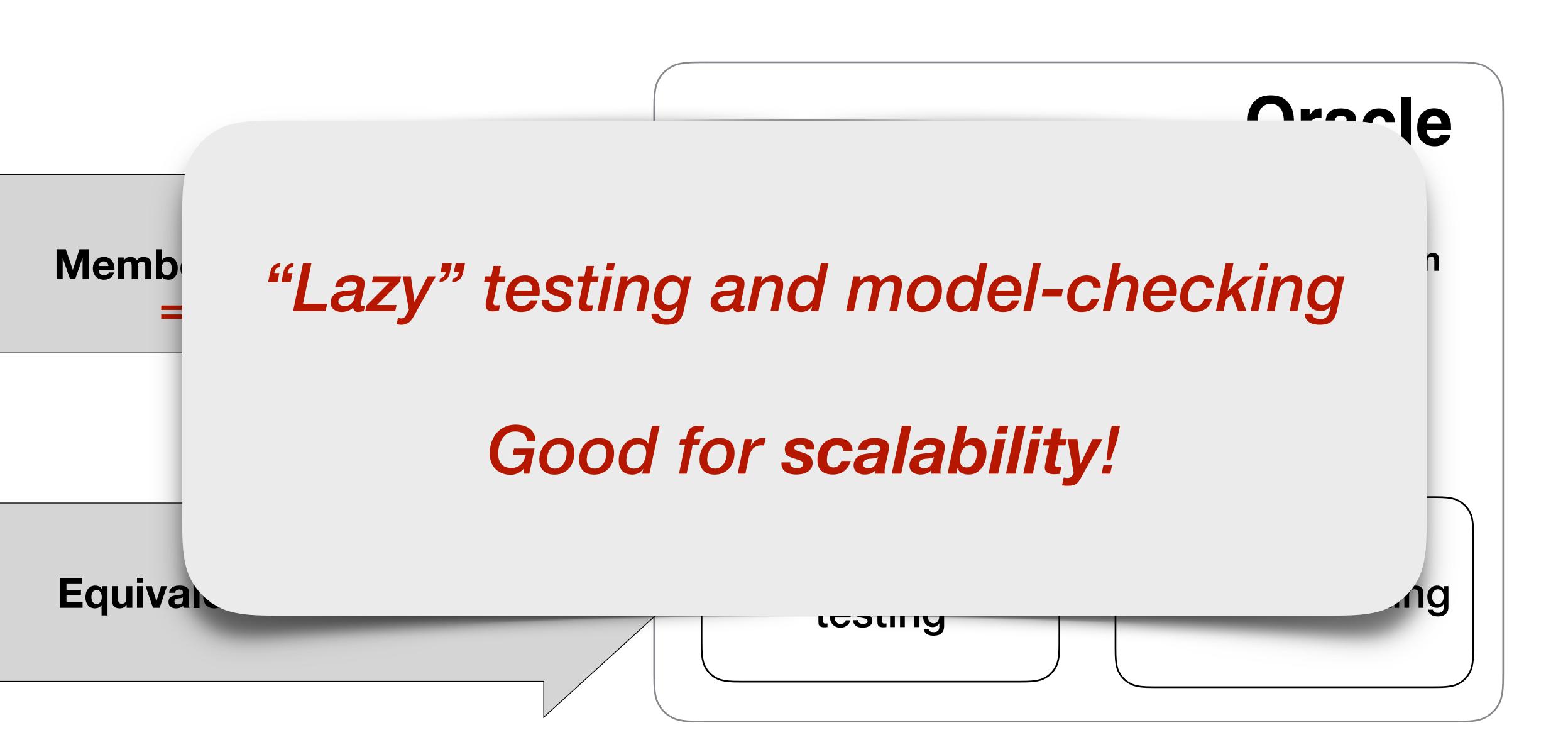
Q: $\mathcal{L}(H) = \mathcal{L}$? where *H* is a hypothesis automaton

A: Y / N + counterexample

In practice...



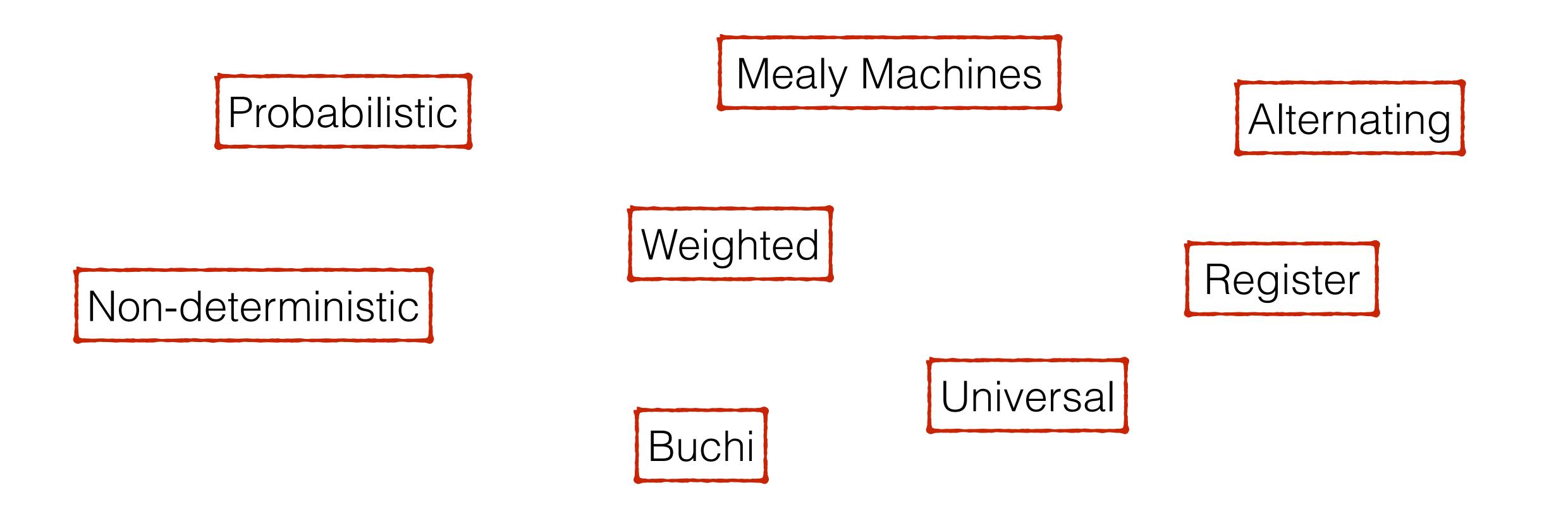
In practice...



Many interesting applications

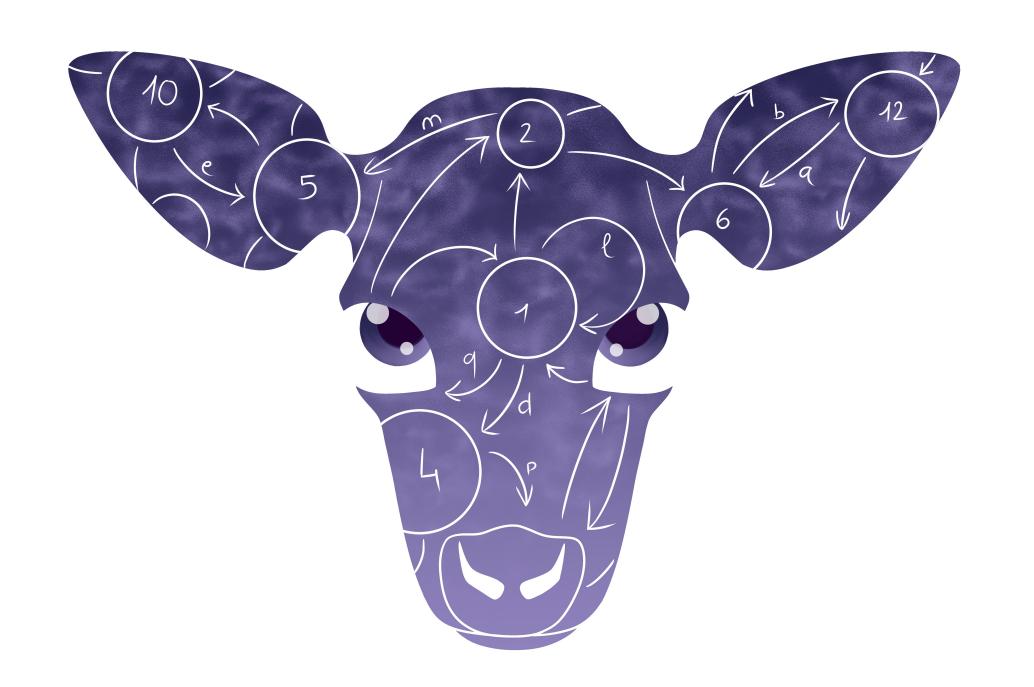
- Detect TLS implementations flaws [USENIX Sec. Sym. '15]
- TCP implementations [CAV '16]
- Analysis of botnet protocols [CCS '10]
- Bank cards ...

To each application domain its model...



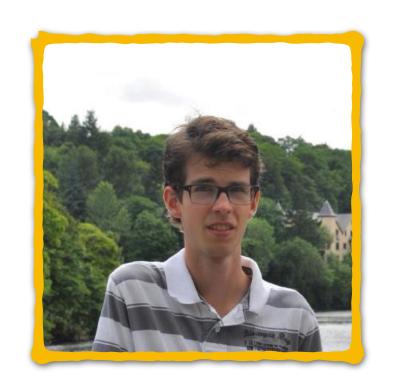
Do I need to write my automata learning algorithm from scratch?

NO! Maths can help!



Categorical Automata Learning Framework

calf-project.org



Gerco van Heerdt UCL











Joshua Moerman **Radboud University**

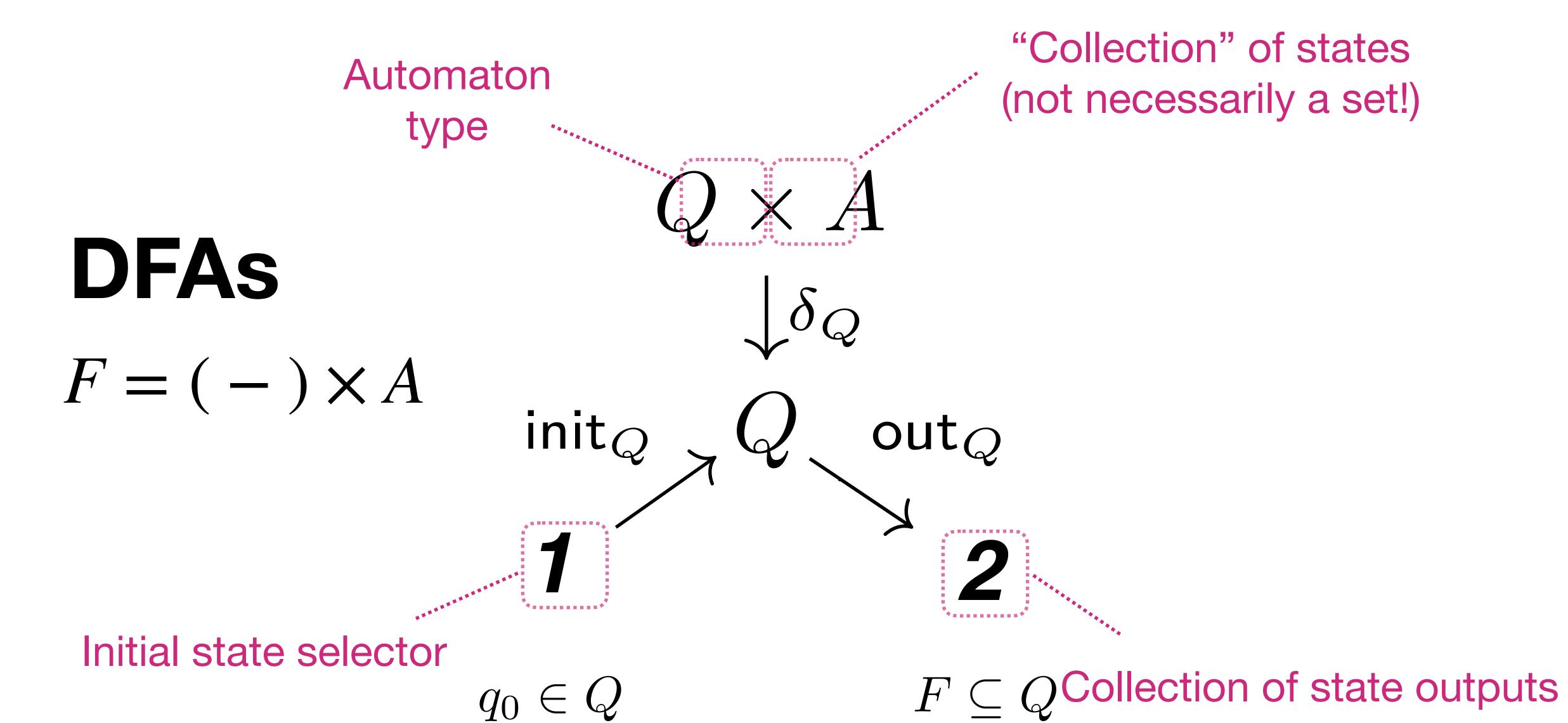
Bartek Klin

Warsaw University Warsaw University

Michal Szynwelski Maverick Chardet **ENS Lyon**

Tiago Ferreira **UCL Intern**

Different automata, same structure

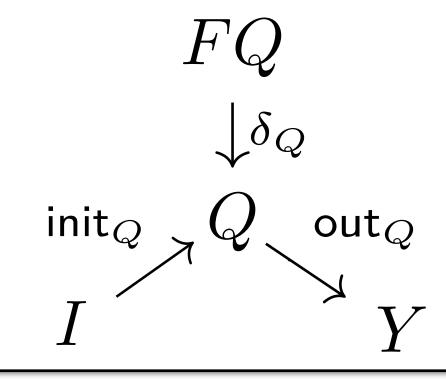


A general framework

Abstract observation data structure

approximates

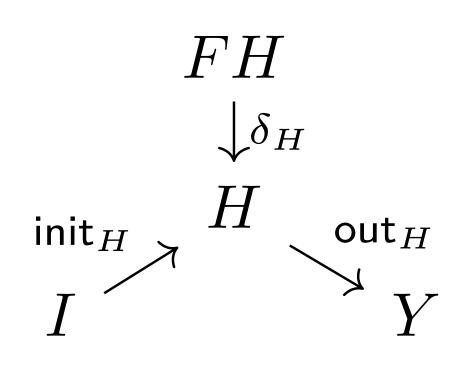
Target minimal automaton

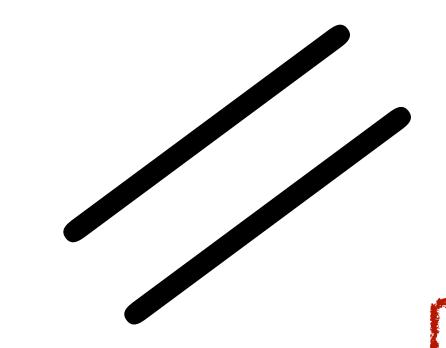




abstract closedness and consistency

Hypothesis automaton





General correctness theorem

Guidelines for implementation

New algorithms!

F. Categorical Automata Learning Framework (CSL '17)

Learning Nominal Automata (POPL '17)

Joshua Moerman, Matteo Sammartino, Alexandra Silva, Bartek Klin, Michal Szynwelski

Nom Nominal automata

Vect Weighted automata

Change main data structure

Observation tables

Discrimination trees

Plug monads in

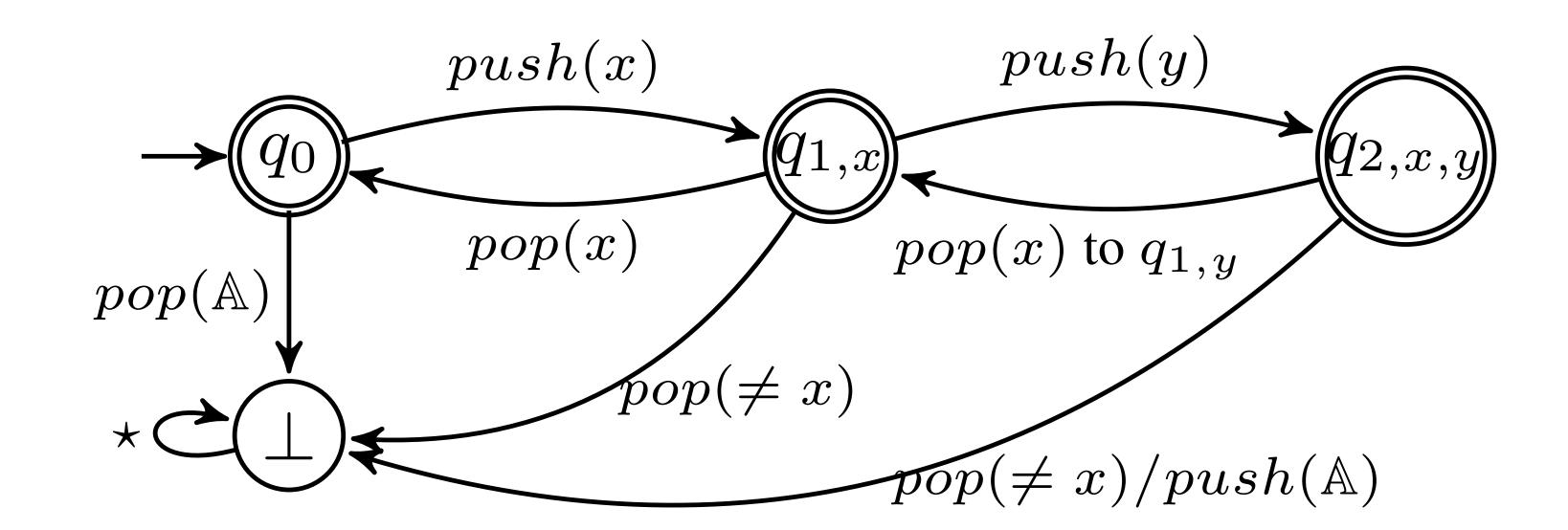
Powerset NFAs

Powerset with intersection Universal automata

Double powerset Alternating automata

Infinite alphabets

infinite-state, but finitely representable automata



Change automaton type

Change main data structure

Set DFAs

Nom Nominal automata

Vect Weighted automata

Observation tables

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Optimising Automata Learning via Monads

Gerco van Heerdt, Matteo Sammartino, Alexandra Silva

(arXiv:1704.08055)

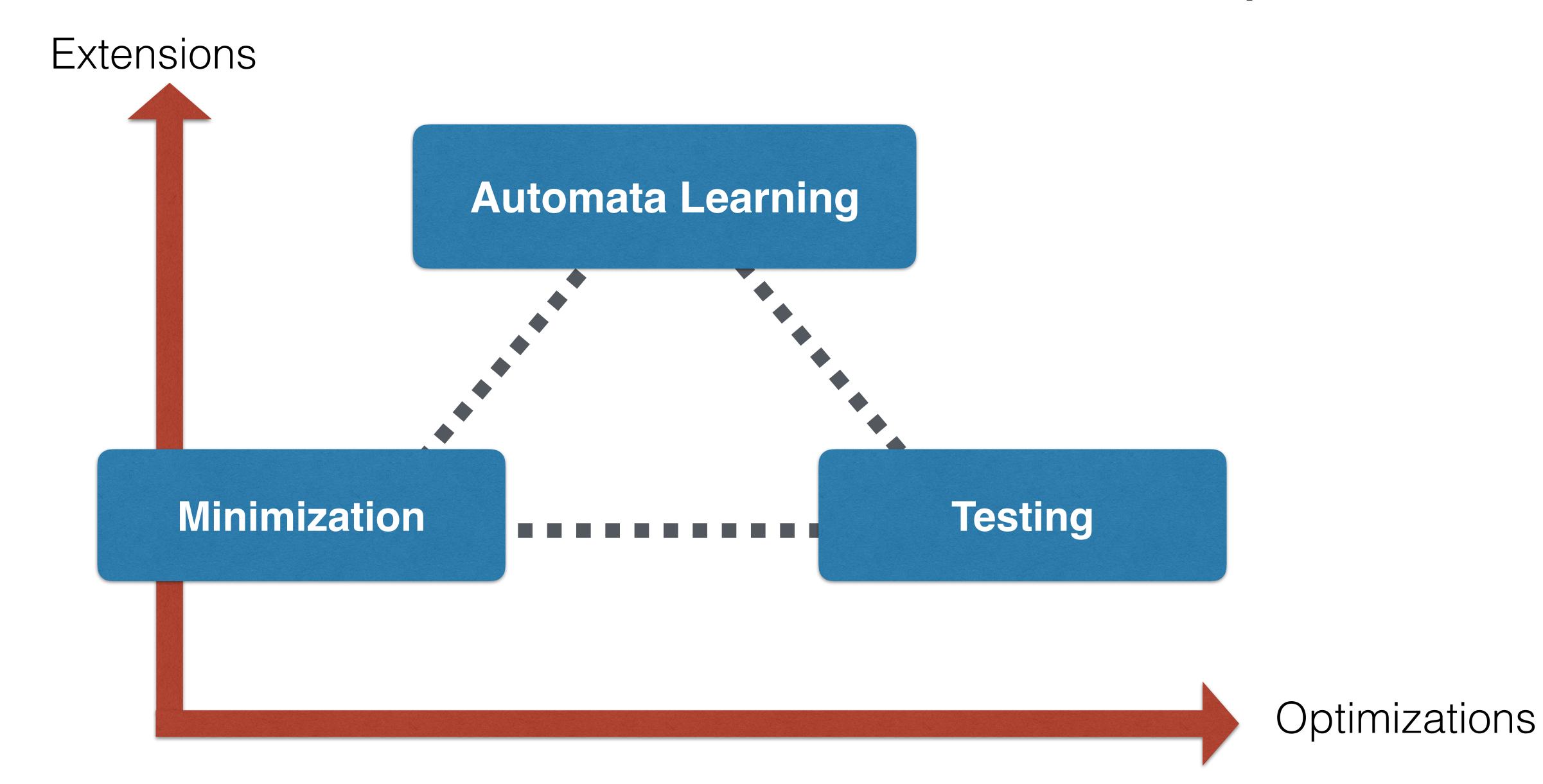
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Connections with other techniques



Basic ingredients

NetKat

```
sw = 6; pt = 8; dst := 10.0.1.5; pt := 5
```

For all packets located at port 8 of switch 6, set the destination address to 10.0.1.5 and forward it out on port 5.

[Hoare et al., JLAMP '11] [Kappe et al., CONCUR '17, ESOP '18]

Thread 1: do **a** and then **b**

Thread 2: do c and then d

Verification using NetKAT

Reachability

► Can host A communicate with host B? Can every host communicate with every other host?

Security

► Does all untrusted traffic pass through the intrusion detection system located at *C*?

Loop detection

Is it possible for a packet to be forwarded around a cycle in the network?

Verification using NetKAT

Soundness and Completeness [Anderson et al. 14]

 $ightharpoonup \vdash p = q \text{ if and only if } \llbracket p \rrbracket = \llbracket q \rrbracket$

Decision Procedure [Foster et al. 15]

- NetKAT coalgebra
- efficient bisimulation-based decision procedure
- implementation in OCaml
- deployed in the Frenetic suite of network management tools

This project

Forwarding/
Filtering behaviour

Concurrency

Large data domains

NetKat CALF CKA

in collaboration with amazon

Other research directions

Software Analysis

- Learning the "correct ways" of using undocumented code
- Learning-based automated test generation

Hardware Analysis

Analysing concurrency in hardware, in collaboration with arm